

How to find the most powerful test?

Given two simple hypotheses  $H_0, H_1$

Design test — dando up possible outcomes

$$x = (x_1, \dots, x_n) \in \mathbb{R}^n = R_0 \cup R_1$$

$x \in R_0 \Rightarrow$  accept  $H_0$  reject  $H_1$

$x \in R_1 \Rightarrow$  accept  $H_1$  reject  $H_0$

Convention: we call  $R_1$  "critical region" =  $C$

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$P(x \in C | H_0) = \alpha$  smaller than some fixed value.  
"type 1 errors"

$1 - \beta = P(x \in C | H_1)$  want as large as possible.

$\beta = P(x \notin C | H_1)$  "type 2 errors"

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It is possible (reasonable) that for values  $x \notin C$   
(discrete)

$$P(X > x | H_0) < P(X = x | H_1)$$

ex:  $H_0 = \text{coin fair}$        $H_1 = P(\text{heads}) = 80\%$

example: 5 flips.       $C = \{5 \text{ heads only}\}$

$$P(\text{Type I error}) = \frac{1}{32}$$

$$P(X=4 | H_0) = \binom{5}{1} \left(\frac{1}{2}\right)^4 \left(\frac{1}{2}\right)^1 = \frac{5}{32} \approx 16\%$$

$$P(X=4 | H_1) = \binom{5}{1} \left(\frac{4}{5}\right)^4 \left(\frac{1}{5}\right)^1 = \frac{5}{5^5} \cdot 4^4 = \frac{256}{625} \\ = \frac{256}{625} \approx 40\%$$

else: if allowed  $C = \{4 \text{ or } 5 \text{ heads}\}$

$$P(\text{Type I}) = P(X=5 \text{ or } 4 | H_0)$$

$$= \frac{1}{32} + \frac{5}{32} = \frac{6}{32} = \frac{3}{16} \approx 19\%$$

It's also totally feasible to  $x \in C$

$$P(X=x | H_0) > P(X=x | H_1)$$

$H_0 = \text{fair coin}$      $H_1 = 99.5\% \text{ heads.}$

example: 10 flips, accept  $H_1$  if 9 or 10 heads.

$$P(X=9, 10 | H_0) \approx 1\%$$

check  $P(X=9 | H_0) \approx 1\%$

$$P(X=9 | H_1) < \frac{1}{2}\%$$

Goal: How to find the most powerful test w/ given  $\alpha$ ?

$$\alpha = P(X \in C | H_0) \quad \begin{array}{l} \text{"size of } C \\ \text{"significance level of the test"} \end{array}$$

Def: Fix a size  $\alpha$  for crit regions, only consider regions of size  $\leq \alpha$ .

We say that  $C$  is a most powerful region for hypotheses  $H_0, H_1$ , if for any other region  $D$  of size  $\leq \alpha$ , we have  $P(X \in C | H_1) \geq P(X \in D | H_1)$

$(1-\beta)$   
"power"

### Neyman-Pearson Lemma

Discrete case

If  $\frac{P(X=x | H_0)}{P(X=x | H_1)}$  is at least as big for  $x \notin C$  as it is for all  $x \in C$  then  $C$  is a most powerful region for  $H_0$

assume  $P(X=x | H_1) \neq 0$  (in region where we accept  $H_0$  want this big)  
 $\forall x \dots \frac{P(X=x | H_0)}{P(X=x | H_1)} \leq \frac{P(X \in C | H_0)}{P(X \in C | H_1)}$  -- small.

more formally

$$\frac{P(X=x|H_0)}{P(X=x|H_1)} \geq \frac{P(X=y|H_0)}{P(X=y|H_1)}$$

all  $x \notin C, y \in C$ .

Proof

set  $K = \inf \left\{ \frac{P(X=x|H_0)}{P(X=x|H_1)} \mid x \notin C \right\}$

above inequality  $\Rightarrow \frac{P(X=y|H_0)}{P(X=y|H_1)} \leq K$

~~$\Rightarrow K \neq 0$~~

$$\frac{P(X=x|H_0)}{P(X=x|H_1)} \geq K$$

$$\Rightarrow \frac{P(X=x|H_0)}{K} \geq P(X=x|H_1) \quad x \notin C$$

$$\frac{P(X=y|H_0)}{K} \leq P(X=y|H_1) \quad y \in C$$

Suppose  $D$  is a region of size  $\leq \alpha$   
 want  $P(X \in C|H_1) \geq P(X \in D|H_1)$

Know:  $P(X \in C | H_0) = \alpha \geq P(X \in D | H_0)$

$$P(X \in C \cap D | H_0) + P(X \in C \cap D^c | H_0)$$

$$= P(X \in C \cap D | H_0) + P(X \in D \cap C^c | H_0)$$

$$\frac{P(X \in C \cap D^c | H_0)}{\geq k} \geq \frac{P(X \in D \cap C^c | H_0)}{k} \geq P(X \in D \cap C^c | H)$$

$$\boxed{P(X \in C \cap D^c | H_1)} \geq \boxed{P(X \in D \cap C^c | H)}$$

$P(X \in C | H_1)$

$$P(X \in C \cap D | H_1) + P(X \in C \cap D^c | H_1)$$

$\geq \geq$

$$P(X \in C \cap D | H) + P(X \in D \cap C^c | H_1)$$

$= P(X \in D | H_1)$

D

## Neyman Pearson (cont case)

heuristically  
replace  $P(X=x | H_0)$  by  $f_0(x)$  pdf  $H_0$   
 $\therefore P(X=x | H_1)$  by  $f_1(x)$  pdf  $H_1$

want to  $\geq \gamma$ :  $\frac{f_0(x)}{f_1(x)} \geq \frac{f_0(y)}{f_1(y)}$   $\Rightarrow C$   
 if  $f_0(x) \neq 0, \dots, x \in C, y \in C$  most power

Actual statement: If  $C$  has pow  $\alpha$ , and

If  $\exists L$  s.t.  $L f_0(x) \geq f_1(x) \quad x \notin C$   
 $L f_0(y) \leq f_1(y) \quad y \in C$

then  $C$  is most power wrt  $\alpha$ .

$$K = \inf \left\{ \frac{f_0(x)}{f_1(x)} \mid x \notin C \right\} \Rightarrow L = \frac{1}{K}$$

Pf in cont case

$$P(x \in C | H_0) = \alpha \geq P(x \in D | H_0)$$

" " "

$$\int_{x \in C} f_0(x) dx = \int_C f_0$$

$$\int_{x \in D} f_0(x) dx \quad " \quad \int_D f_0$$

$$\int \cdot \int \cdot \dots \int f_0(x_1, \dots, x_n) dx_1 \dots dx_n$$

$\prod f_i(x_i)$

$$\int_C f_0 = \alpha \Rightarrow \int_D f_0 \quad "$$

$$\int_{C \cap D} f_0 + \int_{C \cap D^c} f_0 \quad "$$

$$\int_{D \cap C} f_0 + \int_{D \cap C^c} f_0$$

$$\int_{C \cap D^c} f_0 \geq \int_{D \cap C^c} f_0$$

$$L_f = \int_C f_1 = \int_{C \cap D} f_1 + \int_{C \cap D^c} f_1 \Rightarrow \int_{C \cap D} f_1 + L \int_{C \cap D^c} f_0 \geq L \int_{D \cap C^c} f_0 + \int_{C \cap D} f_1$$

$$\geq \int_{D \cap C^c} f_1 + \int_{C \cap D} f_1 = \int_D f_1$$

□