

"Final" exam (aka exam 2)

cheat sheet (1 pg, double sided)

Dec 9, during class (90 minutes)

Focus on material after midterm

Chap 11 (X-modules / chain conditions)

Properties of parts of sub X-modules.

Def A poset (partially ordered set) is a pair (P, \leq) where P is a set, \leq a binary relation s.t.

i) $a \leq a$ all $a \in P$ (reflexive)

ii) $a \leq b, b \leq c \Rightarrow a \leq c$, all $a, b, c \in P$ (transitive)

iii) $a \leq b, b \leq a \Rightarrow a = b$, all $a, b \in P$. (antisymmetric)

Def Notation: $a > b$ means $b \leq a$
 $a < b$ means $a \leq b$ & $a \neq b$
 $a \gg b$ means $a > b$ & $a \neq b$.

Def We say a poset (P, \leq) is totally ordered (or linearly ordered / as a chain) if $\forall a, b \in P$ either $a \leq b$ or $b \leq a$.

Def sub poset... (a subset of P which is also Herby partially ordered)

typical usage: if (P, \leq) a poset

$C \subseteq P$ we typically say C is a chain
(in P)
if C is tot. ordered.

Def we say a poset (P, \leq) satisfies the
ACC (Ascending Chain Condition) if for any
ascending chain of elements

$a_1 \leq a_2 \leq \dots$ we have $a_n = a_{n+1} = a_{n+2} = \dots$
for some n .

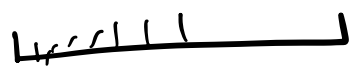
DCC (Descending chain condition) if for any
descending chain of elements

$a_1 \geq a_2 \geq \dots$ we have $a_n = a_{n+1} = a_{n+2} = \dots$
for some n .

(exercise: equivalent to say for ACC:

$\forall C \subseteq P \exists c \in C$ s.t. $c \geq c'$ all $c' \in C$)

ex: $\{1/n \mid n \in \mathbb{Z}_{>0}\} \subset \mathbb{R}$ satisfies ACC not DCC



ex: $\{1 - 1/n \mid n \in \mathbb{Z}_{>0}\} \subset \mathbb{R}$ - - - DCC - - ACC

Note: if P is a poset, define (P^{op}, \leq^{op})
 via $P^{op} = P$ $a \leq^{op} b \Leftrightarrow b \leq a$.

Def We say a poset (P, \leq) satisfies the max'l condition
 if $\nexists \emptyset \neq S \subseteq P$ subset, S has a max'l element

Def We say a poset (P, \leq) satisfies the min'l condition
 if $\nexists \emptyset \neq S \subseteq P$ subset, S has a min'l element

Lemma: for any poset P

i) P satisfies ACC $\Leftrightarrow P$ satisfies max'l condition

ii) P satisfies DCC $\Leftrightarrow P$ satisfies min'l condition

Pr. i) \Leftarrow if P satisfies max'l, $a_1 \leq a_2 \leq \dots$
 an ascending chain.

then let $S = \{a_i \mid i \in \mathbb{Z}_{>0}\}$

then S contains a max'l element an ascending

\Rightarrow if $m \geq n$ then $a_m = a_n$ since $a_m \geq a_n$
 $a_n \geq a_m$ maximality.

conversely, suppose P doesn't satisfy
 max'l condition. will show doesn't satisfy ACC.

let $S \subseteq P$ have no max'l element.

for each $s \in S$, $\{t \in S \mid t > s\} \neq \emptyset$

So $\left(\prod_{s \in S} \{t \in S \mid t > s\} \right) \neq \emptyset$ Axiom of choice.

$$\begin{array}{c} \downarrow \\ f: S \rightarrow S \quad f(s) > s \end{array}$$

Take a sequence
choose any $s \in S$

$s, f(s), f(f(s)), \dots$
contradicts ACC.

~~without Axiom of choice~~

~~by hyp $S \neq \emptyset$~~

~~choose $s \in S$,~~

Def X -modules = Abelian X -group

we say an X -module M is Noetherian if the prot of X -submodules of M satisfies the ACC.

we say an X -module M is Artinian if the prot of X -submodules of M satisfies the DCC.

Thm if M an X -module then M has finite comp. length iff M is both Artinian & Noetherian.

PF: Suppose M is Artinian; Noetherian.

Consider the collection of all χ -submodules of M which have finite length. Suppose M doesn't have f.l.

ACC \Rightarrow can choose $S < M$ max'l finite length submodule. By assumption $S \neq M$

Consider the collection of χ -submodules of M which contain S .

property.

By DCC can find $T < M$ min'l property containing S .

But by corresp. thm, T/S has no proper submods

so is simple \Rightarrow

$$l(T) = l(S) + 1 \Rightarrow T \text{ has finite length.}$$

(a): $H_0 < \dots < H_m = S < T$ is a comp

This contradicts the maximality of S .

\Rightarrow contradict our assumption χ has finite length.

Go to direction left to the reader: (read it)

D.

Theorem Let M be an X -module. Then M is Noth. if and only if every X -submodule $N < M$ is finitely generated.

"Recall" if $S \subset M$ any subset of an X -module M ,
set $\langle S \rangle = \bigcap_{S \subset N < M} N =$ "the submodule generated by S "
 $N < M$ is finitely generated if $\exists S$ finite w/ $N = \langle S \rangle$.

Lemma: If \mathcal{B} is a totally ordered collection of X -submodules of M then $\bigcup_{H \in \mathcal{B}} H$ is a submodule of M .

PF (sketch) if $h_1, h_2 \in \bigcup H \Rightarrow h_1 \in H_1 \in \mathcal{B} \quad h_2 \in H_2 \in \mathcal{B}$
and (say) $H_1 < H_2 \Rightarrow h_1, h_2 \in H_2 \Rightarrow h_1 + h_2 \in H_2 \Rightarrow h_1 + h_2 \in \bigcup H$.
etc.

PF of thm:

if M is Noth wts $\forall N < M, N$ is f.g.

Let $N < M$. let $\mathcal{A} = \{S < N \mid S \text{ f.g.}\}$

So let $S < N$ be a max'l elem. of \mathcal{A} .

Claim: $S = N$. if not, choose $n \in N \setminus S$

$S = \langle R \rangle$ R Art.

consider $S' = \langle R \cup \{n\} \rangle$. is. f.g. but strictly
larger than S contradicts maximality of S

$\Rightarrow S = N$. \checkmark

Conversely:

if all submods $N \subset M$ are f.g. WTS M Noth.

choose $N_1 \subset N_2 \subset \dots$

let $N = \cup N_i$ is a submodule.

but $N = \langle R \rangle$ R Art by hyp.

but for each $r \in R$ $r \in N_i(r)$

let $m \geq i(r)$ $\forall r$ (R Art).

$\Rightarrow R \subset N_m$

$\langle R \rangle \subset N_m$

$\overset{N}{\parallel}$

so $N_m = N_{m+1} = \dots = N$.

\square .

Thm M an X -module, $N \subset M$ X -submod

then 1) M Noth $\Leftrightarrow N \leq M/N$ Noth

2) M Art $\Leftrightarrow N \leq M/N$ Artinian.