

Lecture 16: Vizing's theorem

Thursday, March 17, 2016 12:37 PM

G a graph, no loops, $\Delta = \Delta(G)$ max degree of a vertex of G

Thm (Vizing) $\chi'(G) = \Delta$ or $\Delta + 1$

Pf. want to show $\chi'(G) \leq \Delta + 1$ i.e. G is edge $\Delta + 1$ -colorable.

i.e. can find an edge coloring

$c: E_G \rightarrow \{1, \dots, \Delta + 1\}$ w/ $\Delta + 1$ colors
s.t. it's proper: no adjacent edges get same color.

G by induction on # of edges

Assume that every graph w/ fewer than n edges satisfies Vizing's thm ($\chi'(G) \leq \Delta(G) + 1$)

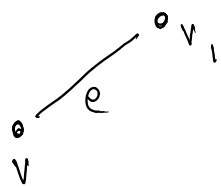
Let G be a graph w/ n edges. want to show $\chi'(G) \leq \Delta(G) + 1$

Strategy, choose $e \in E_G$, consider $G - e$.

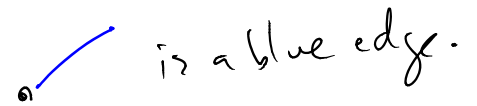
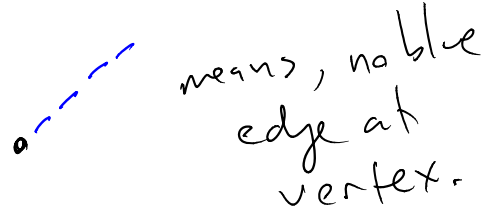
$\Delta(G - e) \leq \Delta(G)$ by induction, $G - e$ is edge $\Delta(G - e) + 1$ colorable

$\Rightarrow G-e$ is edge $\Delta(G)+1$ colorable.

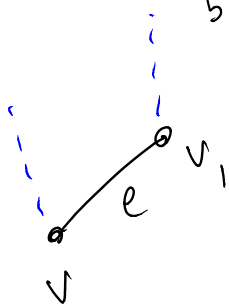
Can pick $c \in E_{G-e} \rightarrow \{1, \dots, \Delta+1\}$ proper edge coloring.



visual notation:

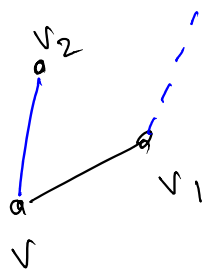


Case: \exists color not appearing at both v_1, v



in this case, can color e blue, done!

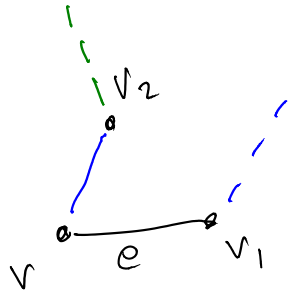
$\Delta+1$ possible colors to assign, each color arises either at v_1 or v .



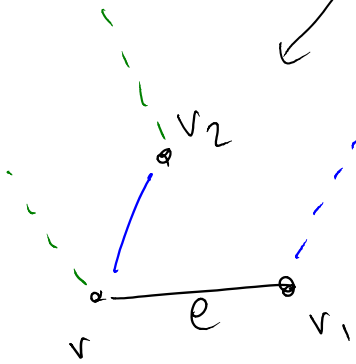
\exists color, which doesn't arise at v_1 and therefore, will have to arise at v .

\exists color which doesn't arise at v_2

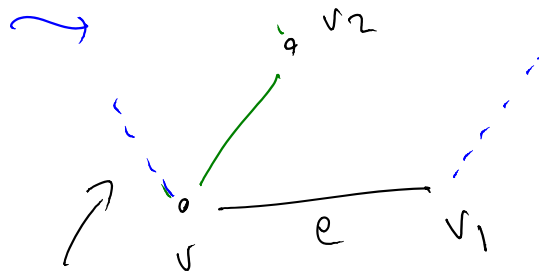
\exists color which doesn't arise at v_2



if green doesn't arise at v



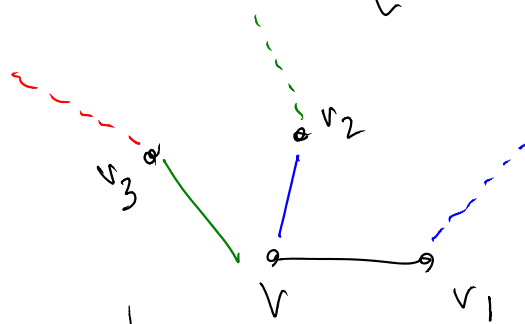
recolor blue edge vv_2 green



just recolored only blue edge.

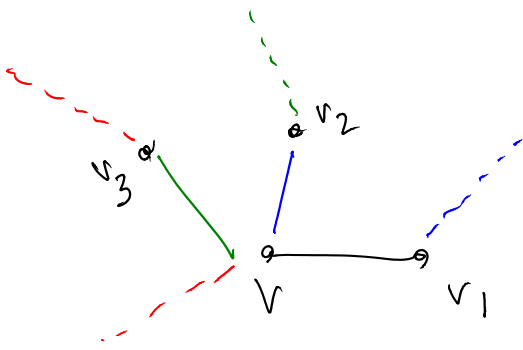
so now, no blue at v .
 \Rightarrow can color e blue \checkmark

or, maybe v has a green edge!



there is a color not at v_3
 suppose it's a new color

maybe v has no red edges...

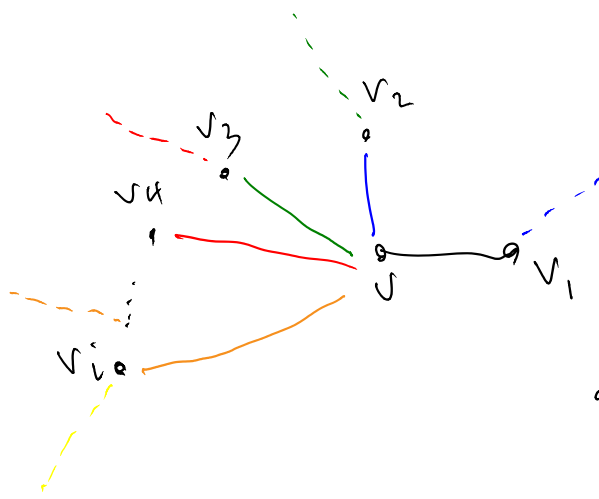


in this case, make
 vv_3 red,
 vv_2 green
 vv_1 blue!

or maybe it does have a red edge.



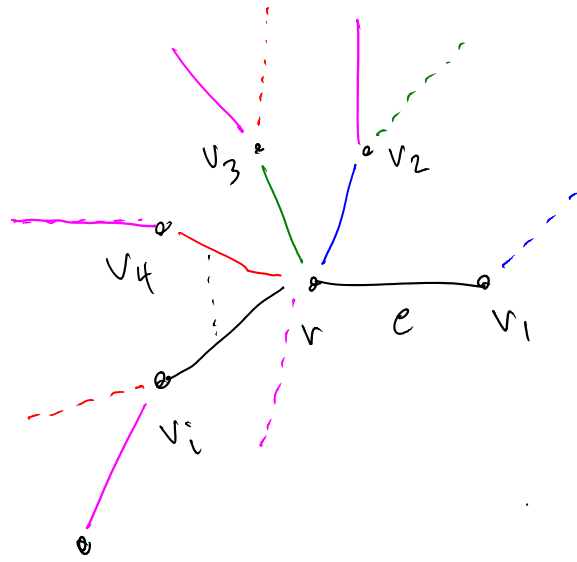
keep going.



pattern continues until
 run out of edges at v
 or • get a repeated color.

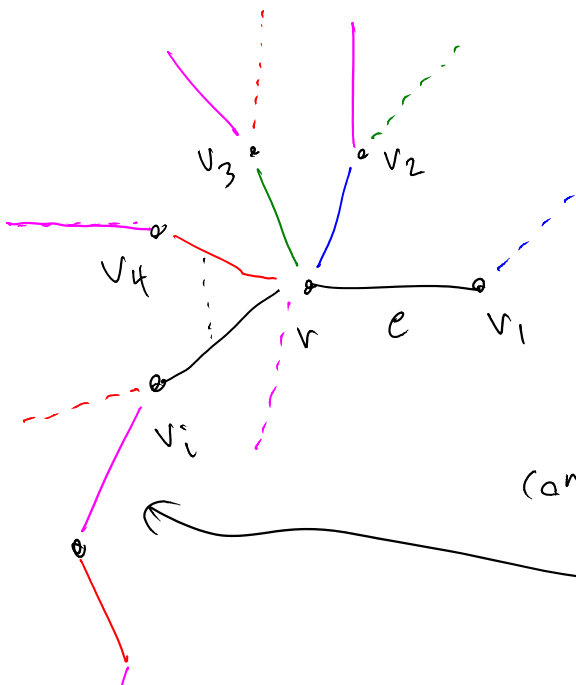
If yellow is a new color i out of edges at v , then
 yellow doesn't arise at v , so can recolor vv_i yellow,
 ... - recolor
 to eventually
 make vv_1

Can assume therefore that eventually, get a repeated color blue.



consider all red & purple edges.
 gives a subgraph where each vertex is \leq deg 1 or 2 (or 0)

= paths, cycles, or isolated vertices.



consider part of subgraph containing

this edge.
 (path)



v