

# Lecture 9: connectivity, part 2

Tuesday, February 9, 2016 12:37 PM

disconnect a graph

$K(G)$

$K'(G)$

independent paths

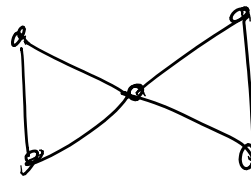


cycles



if ~~simple~~  $G$  has no bridges  
then every edge lies on a cycle.  
or bridge

A vertex  $v \in V_G$  is a cut vertex if  $c(G-v) > c(G)$



Prop: Suppose  $G$  is connected, has at least 3 vertices and no cut vertices, then every pair of vertices lie on a cycle.

Pf: Define distance between vertices = <sup>min</sup> length of a path between them

$\longrightarrow$  dist = 1

Choose vertices  $u, v \in V_G$ , such an distance between them.

dist = 1



remove  $e$ ,  $G-e$  is connected since

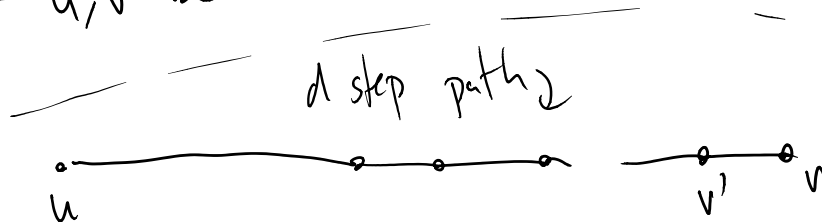
$$k^2(G) \geq k(G) \geq 2$$

$\Leftarrow$   
 $e$  not a bridge.

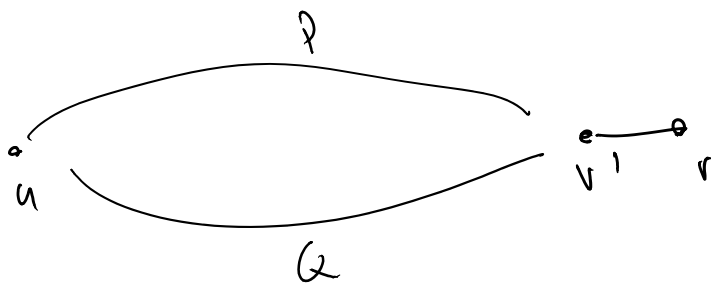
$\Rightarrow e$  on a cycle  $\square$

suppose true for vertices  $\leq d-1$  apart.

let  $u, v$  be vertices distance  $d$ .

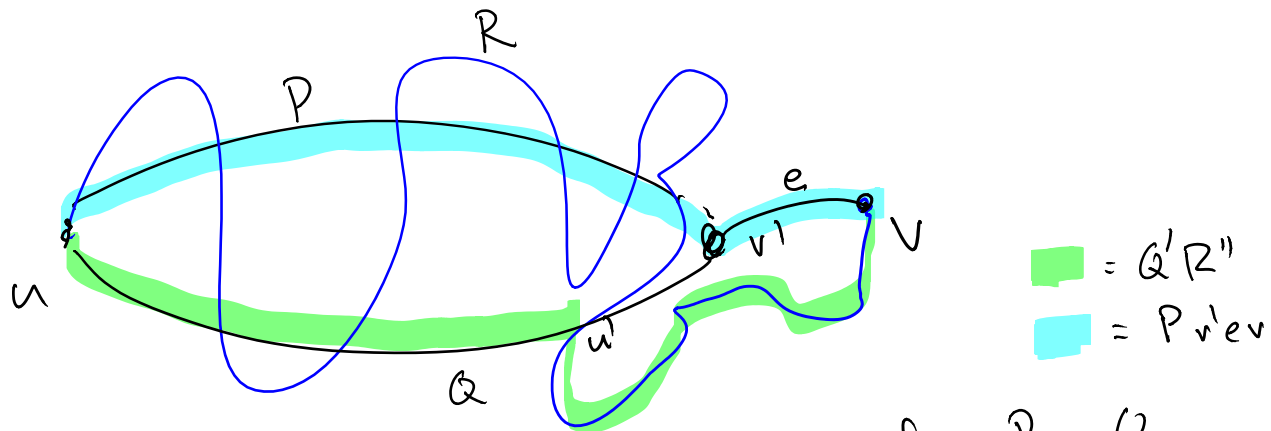


$\Rightarrow u, v'$  are  $d-1$  apart  $\Rightarrow \exists$  disjoint paths  $P, Q$  from  $u$  to  $v'$



now consider  $G-v'$ . This is connected,  $u, v \in V_{G-v'}$

so can find a  $(u, v)$ -path  $\hat{R}$  in  $G - v$



Case 1: if  $R$  happens to be disjoint from  $P$  or  $Q$ ,  
done: 2 paths are  $R$  &  $P v' e v$

(say  $R$  disjoint from  $P$ )

Case 2:  $R$  hits both. suppose in going from  $u$  to  $v$  it hits  $Q$  last at some vertex  $u'$

$R = R' R''$  where  $R'$  is a  $(u, u')$ -path  
 $R''$  is a  $(u', v)$ -path

by construction,  $R''$  is int. disjoint from  $P$  &  $Q$ .  
it's along  $Q$ , so can write  $Q = Q' Q''$

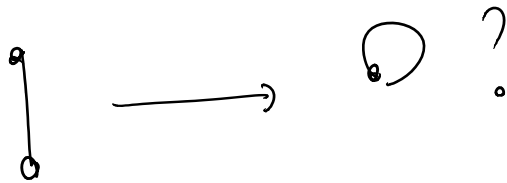
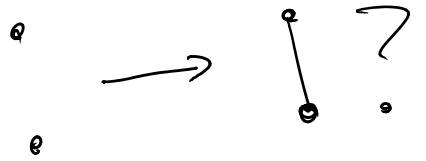
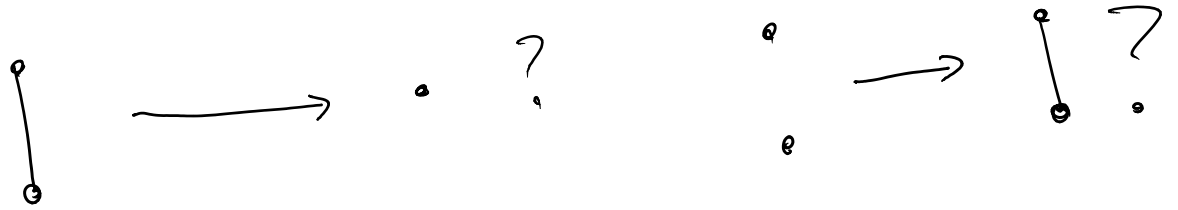
$Q'$  is a  $(u, u')$ -path,  
 $Q''$  is a  $(u', v)$ -path

now, new int. disj. paths are

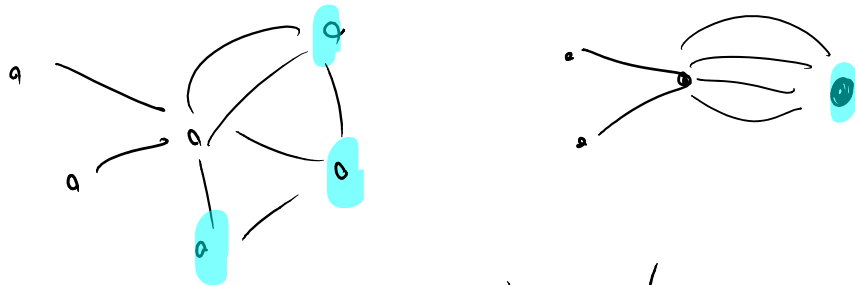
$R''$  □

now, new int. disj. paths in  
 $P$  v'ev  $\frac{1}{2}$ ,  $Q, R$ .  $\square$

maps between graphs (why we don't generally talk about it)



Given a graph  $G$ ,  $A \subset V_G$ , define  $G/A$   
 "shrinking  $A$  to a vertex"



Def  $G/A$  is the graph w/  
 $V_{G/A} = (V_G \setminus A) \cup \{a\}$

$$E_{G/A} = E_G \setminus [A, A]$$

$$\psi_{G/A}(e) = \begin{cases} \psi_G(e) & \text{if } e \in [\bar{A}, \bar{A}] \\ va & \text{if } \psi_G(e) = va \text{ same } a \in A. \end{cases}$$

Thm (Menger)  $G$  is a <sup>connected</sup> graph,  $u, v$  nonadjacent  
then  $\kappa(u, v) = \max \#$  of pairwise disjoint  $(u, v)$ -paths.

Def:  $\kappa(u, v) = \min \#$  of vertices we need to remove so that there are no  $(u, v)$  paths.

(only defined if  $u, v$  nonadjacent)

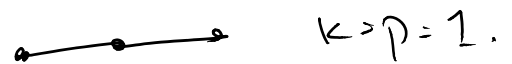
Note:  $\kappa(G) = \min \{ \kappa(u, v) \mid u, v \in V_G \}$  if  $v(G) - 1 \neq \kappa(G)$

Def:  $p(u, v) = \max \#$  mutually disj. paths from  $u$  to  $v$ .

PA WTS  $p = p(u, v) = \kappa(u, v) = \kappa$

Induct on  $e(G)$ .

Base case:  $e(G) = 2 \checkmark$



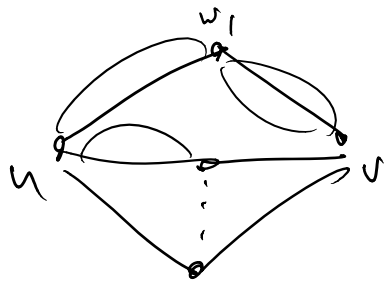
In general Case 1: any edge in  $G$  is incident to either  $v$  or  $u$ .

In general

to either  $u$  or  $v$ .

in this case, every path from  $u$  to  $v$  has length 2

$P_1, P_2, \dots, P_\ell$  indep. paths.



$$P_i = u e_i w_i v$$

$w_i$  all distinct.

$\Rightarrow$  if  $P_i$ 's are a max'l family then  $p = \ell$   
& also  $G - \{w_1, \dots, w_\ell\}$   
has no  $(u, v)$  paths.  
 $\Rightarrow k \leq \ell = p$