

Ab-Cat-Recap

$$\begin{array}{ccc}
 & \swarrow \text{additive!} & \searrow \text{Fix } Y \rightarrow F(Y) \\
 A & \xrightarrow{\quad F \quad} & \xrightarrow{e_Y} Ab \\
 A & \xleftarrow{h_A} & (X \mapsto \text{Hom}(X, A))
 \end{array}$$

Composition:

$$\begin{array}{ccc}
 A & \xrightarrow{\quad Ab \quad} & \text{Hom}(Y, -) \text{ left, but not right exact} \\
 A & \xleftarrow{\quad \text{Hom}(Y, A) \quad} &
 \end{array}$$

Exercise show that e_Y exact \Leftrightarrow a map is monic/epic in

$$\begin{aligned}
 \text{Fun}(A^{\text{op}}, Ab) &\Leftrightarrow \\
 e_Y(\text{map}) \text{ is monic/epic} &\text{ for all } Y.
 \end{aligned}$$

Define $\mathcal{A}^l = \text{full subcat. of } \text{Fun}(A^{\text{op}}, Ab) \text{ whose objects}$

are $\{F(\forall X \in A, \forall a \in FX, \exists Y \xrightarrow{f} X \text{ s.t. } (Ff)(a) = 0\}$

"weakly cocomplete functors"

Fact: \mathcal{A}^l is a Serre subcategory \Leftrightarrow

$A \rightarrow \text{Fun}(A^{\text{op}}, \text{Ab})/\text{el}$ is exact.

Moreover:

If we let $\mathcal{L} \subset \text{Fun}(A^{\text{op}}, \text{Ab})$ be the full subcategory of left exact functors then

$$\begin{array}{ccc} \mathcal{L} & \longrightarrow & \text{Fun}(A^{\text{op}}, \text{Ab}) \\ & \xrightarrow{\sim} & \xrightarrow{\text{equivalence}} \text{Fun}(A^{\text{op}}, \text{Ab})/\text{el} \end{array}$$

Why?

Quick outline of localization / quotient of Abelian cats.

A Abelian category el a full subcategory

el is called a Serre subcat (aka. thick, épaisse)

if $\# 0 \rightarrow A' \rightarrow A \rightarrow A'' \rightarrow 0$ in SES_A ,

$A', A'' \in \text{el} \iff A \in \text{el}$.

Thm (Some) \exists An Abelian cat A/cl & an exact functor $A \xrightarrow{T} A/\text{cl}$ s.t.

If $a \in \text{cl}$, $Ta \cong 0$ & if $A \xrightarrow{T'} B$ is any other exact functor w/ $T'a \cong 0$ for all $a \in \text{cl}$, then $\exists!$ $F: A/\text{cl} \rightarrow B$ s.t.

$$\begin{array}{ccc} A & \xrightarrow{T} & A/\text{cl} \\ T \downarrow & & \downarrow F \\ B & & \end{array}$$

Construction of A/cl :

$$\text{ob}(A/\text{cl}) = \text{ob}(A)$$

given $A' \xrightarrow{f} A$, $B \xrightarrow{g} B'$ in A , have a natural map

$$\text{Hom}_A(A, B) \rightarrow \text{Hom}_{A'}(A', \text{coker } g)$$

$$A' \xrightarrow{f} A \rightarrow B \xrightarrow{\text{coker } g}$$

let $\Lambda_{A,B} = \left\{ (A' \rightarrow A, B' \rightarrow B) \text{ monic s.t. } A/A', B' \in \mathcal{E} \right\}$

this is a directed set and we define

$$\text{Hom}_{A/\mathcal{E}}(A, B) = \varinjlim_{(A', B') \in \Lambda_{A,B}} \text{Hom}_A(A', B/\text{im } B')$$

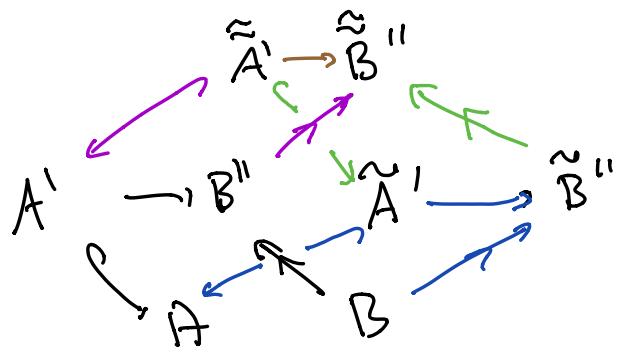
Notation: a map $A \rightarrow B$ in \mathcal{A} is \mathcal{E} -epic if
coker \mathcal{E}

is \mathcal{E} monic if ker \mathcal{E}

{, is \mathcal{E} -iso if both the above.

homom $A \rightarrow B$ in \mathcal{A}/\mathcal{E} are equivalence classes of

$$\begin{matrix} \text{hom} & A' & \longrightarrow & B'' \\ \text{d-epic} & \downarrow & & \uparrow \text{d-monotone} \\ A & & B & \end{matrix}$$



Example R comm. ring, $S \subset R$ a multiplicative set.

$$\text{cl} = \left\langle \{M \in \text{Mod}_R \mid \forall m \in M, sm = 0 \text{ some } s \in S\} \right\rangle$$

then cl is thick \nparallel , $\text{Mod}_R/\text{cl} \cong \text{Mod}_{RS^{-1}}$

In this case the quotient cat is a "localizer"

in fact, there is a right adjoint

$$\text{Mod}_{RS^{-1}} \rightarrow \text{Mod}_R \text{ and}$$

$$\begin{array}{ccc} \text{Mod}_{RS^{-1}} & \xrightarrow{\perp} & \text{Mod}_R \text{ "loc"} \\ & \searrow & \downarrow \\ & & \text{Mod}_{RS^{-1}} \end{array}$$

$\sim id$

$$\text{Mod}_R \xrightarrow{\text{loc}} \text{Mod}_{R\mathcal{S}^{-1}} \xrightarrow{\cong} \text{Mod}_R$$

"localization"

Prop if $T : A \rightarrow A/\text{el}$ has a right adjoint S

then $TS \simeq \text{id}_{A/\text{el}}$ and so $\underbrace{S(A/\text{el})}_{B} \simeq A/\text{el}$

Language: we say in this case that T ($\underset{\text{or } B}{\text{or el}}$) is localizing.

Back to homological algebra

"Last time" we constructed left derived functors
for ^{right exact} $\text{functs } F: \mathcal{A} \rightarrow \mathcal{B}$ in the case that
 \mathcal{A} had enough projectives.

i.e. a \mathcal{S} -funct., universal $\{F_i\}$, $F_0 = F$

Reversing arrows, we find if $F: \mathcal{A} \rightarrow \mathcal{B}$ is
left exact, and if \mathcal{A} has enough injectives
then \exists univ. \mathcal{S} -functs $\{F^i\}$ $F^\circ = F$

"right derived functors of F "
 $R^i F = F^i$

We showed that Mod_R has enough projectives.

In fact Mod_R also has enough injectives.